clMAGMA: High Performance Dense Linear Algebra with OpenCL

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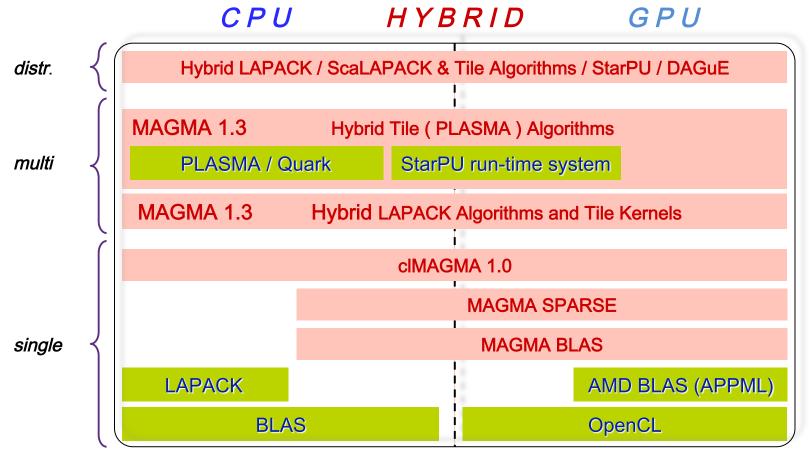
1st International Workshop on OpenCL (IWOCL), May 13—14, 2013 Georgia Institute of Technology, Atlanta, GA

© Outline

- Methodology overview
 - Hybridization of Linear Algebra Algorithms
 - Use both GPUs and multicore CPUs
- clMAGMA
 - OpenCL port of MAGMA
 - Performance results
 - Challenges and future directions
- Conclusions



clMAGMA Software Stack

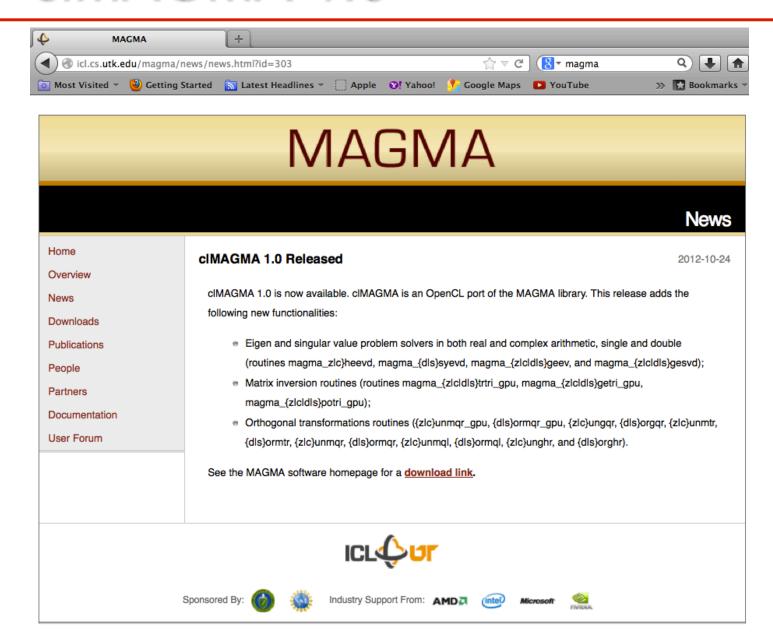


Linux, Windows, Mac OS X | C/C++, Fortran | Matlab, Python

[AMD APPML -- Accelerated Parallel Processing Math Libraries http://developer.amd.com/libraries/appmathlibs/



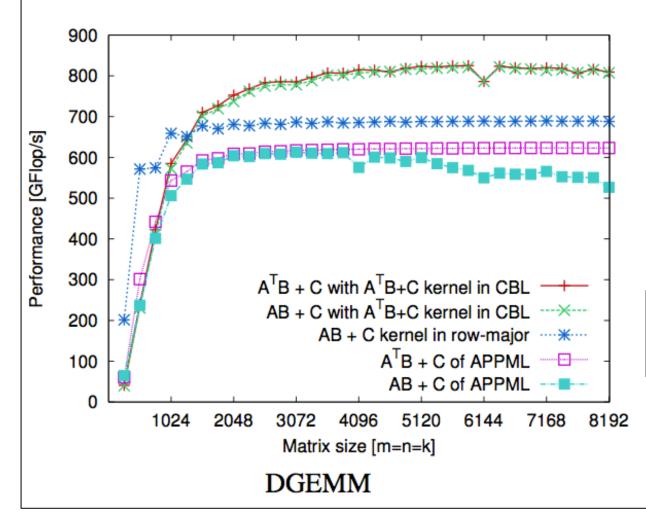
cIMAGMA 1.0





DGEMM in OpenCL

Kazuya Matsumoto, Naohito Nakasato, Stanislav G.Sedukhin, *Implementing a Code Generator for Fast Matrix Multiplication in OpenCL on the GPU*, University of Aizu, Japan, July 2, 2012.



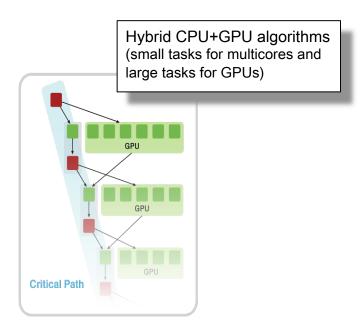
GPU: Tahiti (AMD Radeon HD 7900) 264 GB/s memory bandwidth 3.79 Tflop/s SP, 947 Gflop/s DP 32 x 64 (2048 stream proc.)



MAGMA Methodology

A methodology to use all available resources:

- MAGMA uses HYBRIDIZATION methodology based on
 - Representing linear algebra algorithms as collections of TASKS and DATA DEPENDENCIES among them
 - Properly SCHEDULING tasks' execution over multicore and GPU hardware components
- Successfully applied to fundamental linear algebra algorithms
 - One and two-sided factorizations and solvers
 - Iterative linear and eigen-solvers
- Productivity
 - 1) High-level; 2) Leveraging prior developments; 3) Exceeding in performance homogeneous solutions

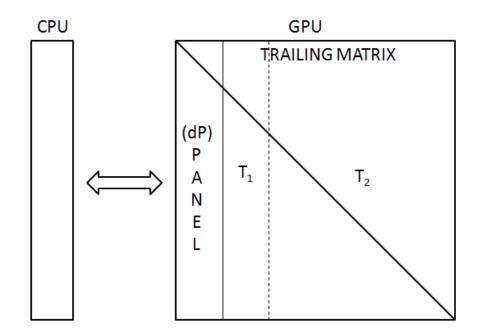




Hybrid Algorithms

One-sided factorizations (LU, QR, Cholesky)

- Hybridization
 - Panels (Level 2 BLAS) are factored on CPU using LAPACK
 - Trailing matrix updates (Level 3 BLAS) are done on the GPU using "look-ahead"



A Hybrid Algorithm Example

Left-looking hybrid Cholesky factorization in clMAGMA

```
for (j=0; j< n; j+= nb) {
 2
          ib = min(nb, n - i);
 3
          magma zherk( MagmaUpper, MagmaConjTrans, jb, j, m one, dA(0, j), ldda, one, dA(j, j), ldda, queue );
 4
          magma zgetmatrix async( jb, jb, dA(j,j), ldda, work, 0, jb, queue, &event );
 5
          if (j+jb < n)
 6
              magma zgemm( MagmaConjTrans, MagmaNoTrans, jb, n-j-jb, j, mz one,
 7
                               dA(0, j), Idda, dA(0, j+jb), Idda, z one, dA(j, j+jb), Idda, queue);
 8
          magma event sync( event );
          lapackf77 zpotrf( MagmaUpperStr, &ib, work, &ib, info );
 9
          if ( *info != 0 )
10
11
              *info += i:
          magma zsetmatrix async(jb, jb, work, 0, jb, dA(j,j), ldda, queue, &event);
12
13
          if (i+ib < n)
12
              magma event sync( event );
13
              magma ztrsm( MagmaLeft, MagmaUpper, MagmaConjTrans, MagmaNonUnit,
14
                             jb, n-j-jb, z one, dA(j, j), Idda, dA(j, j+jb), Idda, queue );
15
16
```

- The difference with LAPACK the 4 additional lines in red
- •Line 9 (done on CPU) is overlapped with work on the GPU (from line 6)



Programming model

Host program

```
for (j=0; j< n; j += nb) {
    jb = min(nb, n - j);
    magma zherk( MagmaUpper, MagmaConjTrans,
                   jb, j, m one, dA(0, j), ldda, one, dA(j, j), ldda, queue );
    magma zgetmatrix async( jb, jb, dA(j,j), ldda, work, 0, jb, queue, &event );
    if (j+jb < n)
       magma zgemm( MagmaConjTrans, MagmaNoTrans, jb, n-j-jb, j, mz one,
                        dA(0, j), Idda, dA(0, j+jb), Idda, z one, dA(j, j+jb), Idda, queue);
    magma event sync( event );
    lapackf77 zpotrf( MagmaUpperStr, &jb, work, &jb, info );
    if ( *info != 0 )
        *info += i:
    magma zsetmatrix async( jb, jb, work, 0, jb, dA(j,j), ldda, queue, &event );
    if (i+ib < n)
       magma event sync( event );
       magma ztrsm( MagmaLeft, MagmaUpper, MagmaConjTrans, MagmaNonUnit,
                       jb, n-j-jb, z one, dA(j, j), ldda, dA(j, j+jb), ldda, queue );
```

OpenCL interface – communications

```
magma err t
magma_zgetmatrix_async(
    magma_int_t m, magma_int_t n,
   magmaDoubleComplex const_ptr dA_src, size t dA_offset, magma_int_
   magmaDoubleComplex*
                               hA_dst, size_t hA_offset, magma_int_
   magma_queue_t queue, magma_event_t *event )
   size_t buffer_origin[3] = { dA_offset*sizeof(magmaDoubleComplex),
   size_t host_orig[3]
                          = { 0, 0, 0 };
   size_t region[3]
                          = { m*sizeof(magmaDoubleComplex), n, 1 };
   cl_int err = clEnqueueReadBufferRect(
       queue, dA_src, CL_FALSE, // non-blocking
       buffer_origin, host_orig, region,
       ldda*sizeof(magmaDoubleComplex), 0,
       ldha*sizeof(magmaDoubleComplex), 0,
       hA_dst, 0, NULL, event );
 OpenCL interface - AMD APPML BLAS
magma_zherk(
    magma_uplo_t uplo, magma_trans_t trans,
    magma_int_t n, magma_int_t k,
    double alpha, magmaDoubleComplex_const_ptr dA, size_t dA_offset,
    double beta, magmaDoubleComplex_ptr
                                             dC, size_t dC_offset,
    magma_queue_t queue )
    cl_int err = clAmdBlasZherk(
```

clAmdBlasColumnMajor,

n, k,

return err;

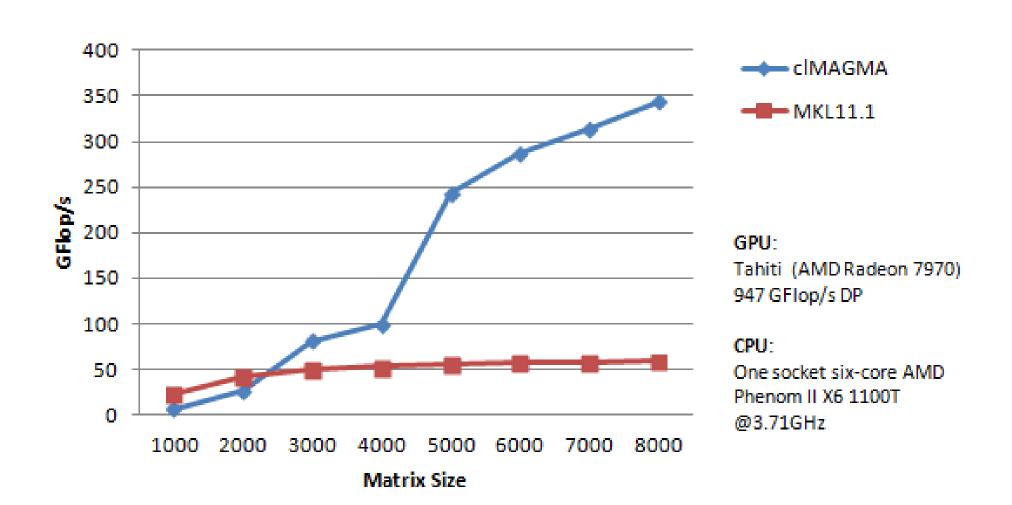
amdblas_uplo_const(uplo),

alpha, dA, dA_offset, lda,
beta, dC, dC_offset, ldc,
1, &queue, 0, NULL, NULL);

amdblas_trans_const(trans),

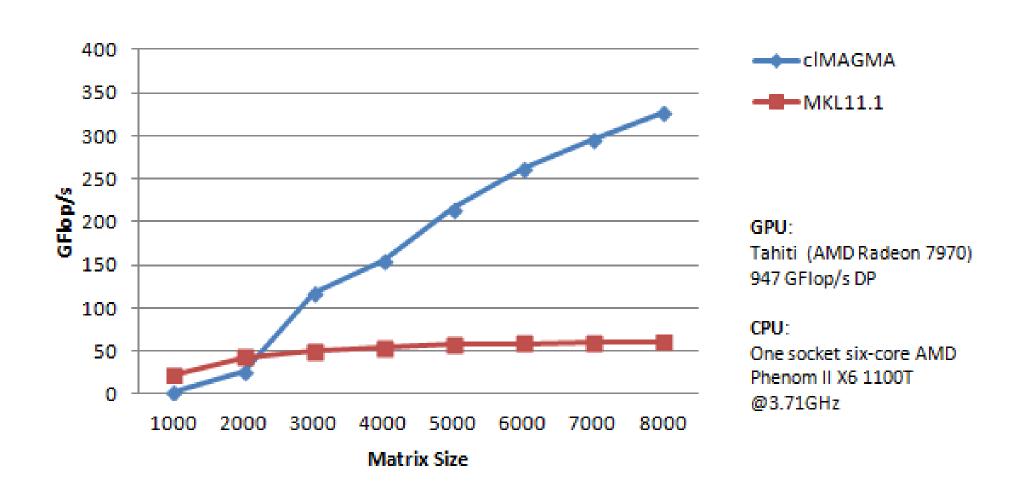


Performance of clMAGMA Cholesky Factorization in double precision



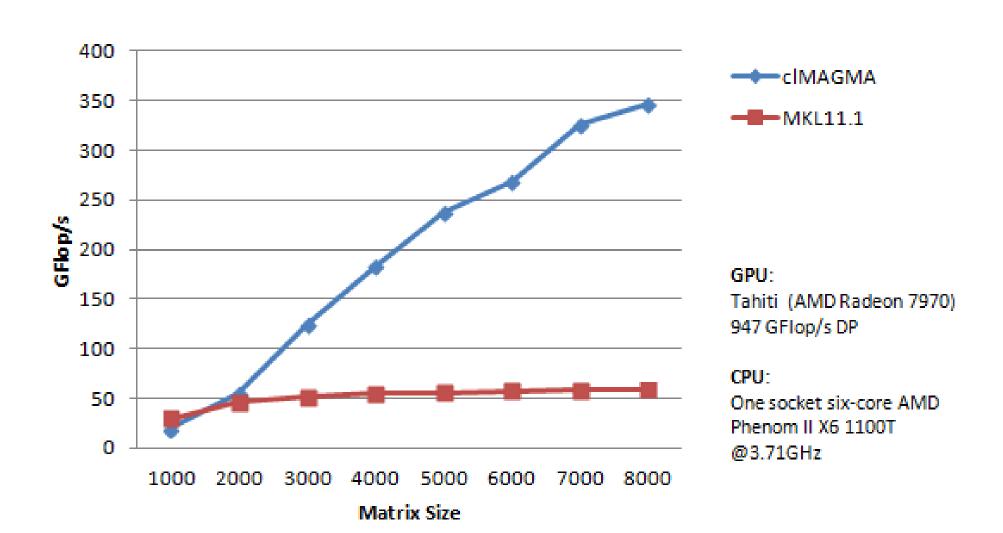


Performance of clMAGMA LU Factorization in double precision





Performance of clMAGMA QR Factorization in double precision

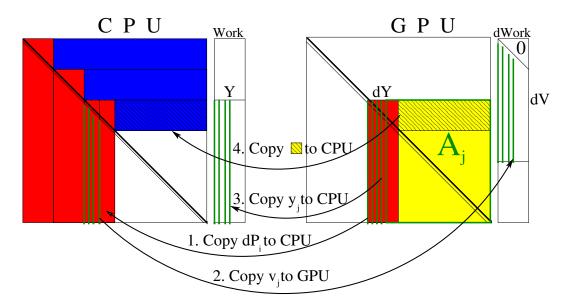




Hybrid Algorithms

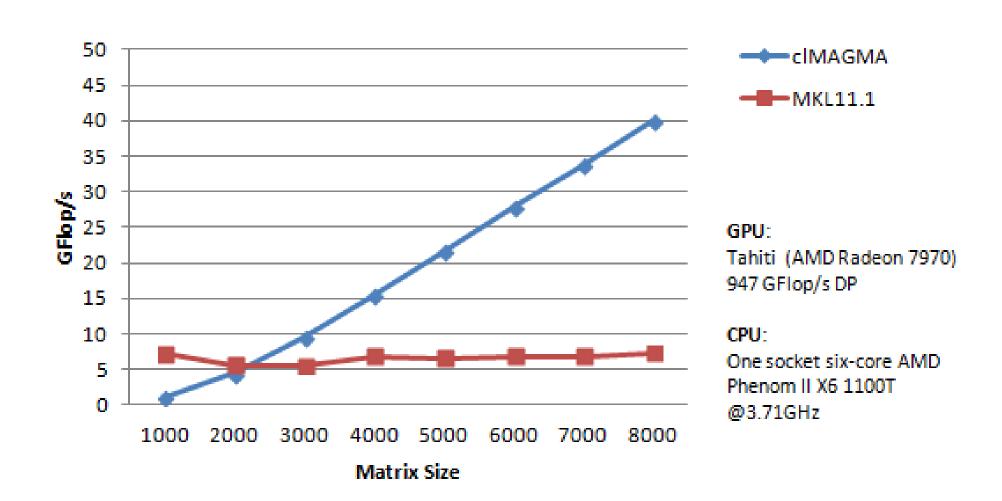
Two-sided factorizations (Hessenberg, bi-, and tridiagonalization)

- **Hybridization**
 - Panels (Level 2 BLAS) are also hybrid, using both CPU & GPU (vs. just CPU as in the one-sided factorizations)
 - Trailing matrix updates (Level 3 BLAS) are done on the GPU using "look-ahead"





Performance of clMAGMA Hessenberg Factorization in double precision



Current work Dynamic Scheduling

Conceptually similar to out-of-order processor scheduling because it has:

- Dynamic runtime DAG scheduler
- Out-of-order execution flow of fine-grained tasks
- Task scheduling as soon as dependencies are satisfied
- -Producer-Consumer

Data Flow Programming Model

- –The DAG approach
- -Scheduling is data driven
- Inherently parallel

Current Work High Level of Productivity

From Sequential Nested-Loop Code to Parallel Execution:

```
for (k = 0; k < min(MT, NT); k++){
          zgeqrt(A[k;k], ...);
          for (n = k+1; n < NT; n++)
                zunmqr(A[k;k], A[k;n], ...);
          for (m = k+1; m < MT; m++){
                ztsqrt(A[k;k],,A[m;k], ...);
                for (n = k+1; n < NT; n++)
                         ztsmqr(A[m;k], A[k;n], A[m;n], ...);
```

Current Work High Level of Productivity

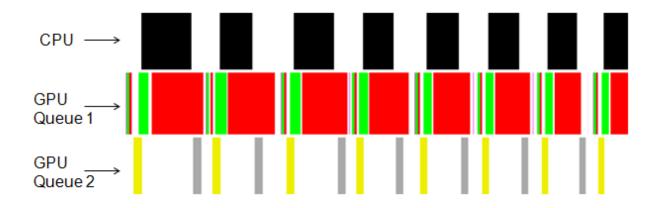
From Sequential Nested-Loop Code to Parallel Execution:

```
for (k = 0; k < min(MT, NT); k++){
          Insert_Task(&zgeqrt, k , k, ...);
         for (n = k+1; n < NT; n++)
                Insert_Task(&zunmqr, k, n, ...);
         for (m = k+1; m < MT; m++){
                Insert_Task(&ztsqrt, m, k, ...);
                for (n = k+1; n < NT; n++)
                         Insert_Task(&ztsmqr, m, n, k, ...);
```

Current work Performance optimizations

 Overlap CPU work, GPU work, and **CPU-GPU communications**

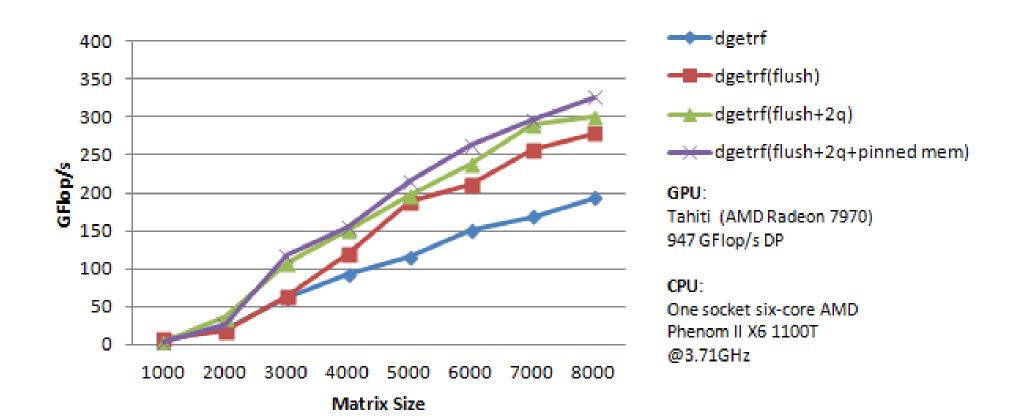
A dgetrf trace example



Gray / Yellow – dset / dget Black – panel Red – dgemm Green – dtrsm



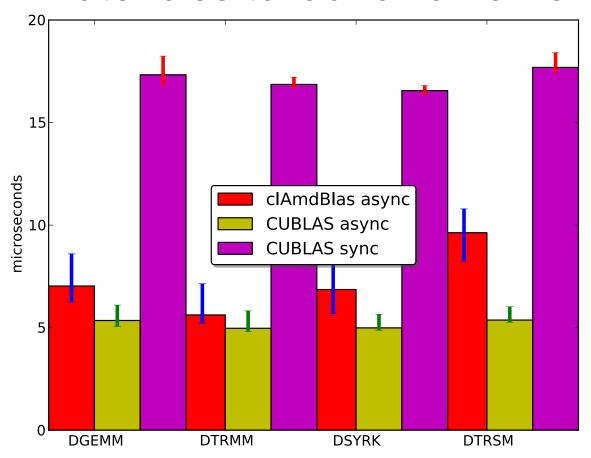
Performance optimizations in LU



OpenCL-specific optimizations

Benchmarks to discover OpenCL specifics

Latencies to launch a kernel





Panels entirely on GPU?

- Important to have for both dense and certain sparse linear system and eigen-problem solvers
- Can we factor panels faster on GPU as panels are memory bound?
- Latencies may be a bottleneck
 - e.g., 64 columns panel would require the invocation of ~400 kernels

Performance of QR panels in double precision on Kepler (in CUDA), Tahiti (in OpenCL), and 16 Intel Sandy Bridge cores

M	N	CUDA Time (ms)	OpenCL Time (ms)	16 Sandy Bridge (ms)
1,000	64	5	94	9
10,000	64	7	104	17
100,000	64	36	131	89
1,000,000	64	365	528	1,431

Difference is due to latencies (in our software/hardware configuration) as shown by increasing the problem size.



Summary and Future Directions

- A hybrid methodology and its application to DLA using OpenCL
- cIMAGMA: LAPACK for heterogeneous computing
 - Achieving high-performance linear algebra using OpenCL
 - clMAGMA 1.0 includes the main
 - one- and two-sided factorizations
 - orthogonal transformation routines
 - linear and eigen-problem solvers
- What is next?
 - Further performance/efficiency improvements
 - MultiGPU and distributed environments



Collaborators / Support

- MAGMA [Matrix Algebra on GPU and Multicore Architectures] team http://icl.cs.utk.edu/magma/
- PLASMA [Parallel Linear Algebra for Scalable Multicore Architectures] team http://icl.cs.utk.edu/plasma
- Collaborating Partners
 University of Tennessee, Knoxville
 University of California, Berkeley
 University of Colorado, Denver

INRIA, France KAUST, Saudi Arabia







