# Parallelization of the Shortest Path Graph Kernel on the GPU 

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## Outline

- Introduction
- Graph
- Graph similarity and graph kernel
- Shortest Path Graph Kernel
- Parallelization on GPU and CPU
- Four GPU implementations
- One OpenMP implementation
- Experiments results
- Synthetic datasets
- Scientific datasets
- Conclusion and Future Work


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- Shortest Path Graph Kernel Darallalization on GDII and CDI - Four GPU implementations - One OpenMP implementation Fvneriments recults
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$\qquad$
Conclusion and Future Work


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## Graph

- A graph $G$ is a set of vertices $\boldsymbol{V}$ and edges $\boldsymbol{E}$, where $\boldsymbol{E} \subset \mathbf{V}^{2}$
- A graph G may have labels on vertices and/or edges
- The adjacency matrix $\boldsymbol{A}$ of G is defined as

$$
\left[\boldsymbol{A}_{i j}\right]=\left\{\begin{array}{lc}
1 & \text { if }\left(v_{i}, v_{j}\right) \in \boldsymbol{E} \\
0 & \text { otherwise }
\end{array}\right.
$$

## Labelled Undirected Graphs

vertices
(1)
(2)

(4)

## Labelled Undirected Graphs

vertices
edges

(4)


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## Labelled Undirected Graphs

vertices
labels


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## Labelled Undirected Graphs

vertices
(1)
(2)
(4)


$$
\boldsymbol{A}=\left[\begin{array}{llll}
0 & 1 & 1 & 1 \\
1 & 0 & 0 & 0 \\
1 & 0 & 0 & 0 \\
1 & 0 & 0 & 0
\end{array}\right]
$$

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## Graph Similarity

- How similar are two graphs?
- For Machine Learning problems like clustering and classification on graphs, graph similarity is crucial.
- Applications
- Protein function prediction
- Drug screening
- Documents classification (Junk mail?)
- Image classification
- Cyber security
- Challenges
- Graph isomorphism is NP-complete
- Graph comparison via isomorphism is prohibitively expensive


## Graph Kernel

- To Calculate the similarities between two graphs in polynomial time
- Random Walk Kernel
- Compare all walks in two graphs $\boldsymbol{G}$ and $\boldsymbol{G}^{\prime}$
- Shortest Path Kernel
- Compare all pairs shortest paths for $\boldsymbol{G}$ and $\boldsymbol{G}^{\prime}$ via Floyd-Warshall
- Subtree Kernel
- Compare subtree-like patterns in two graphs $\boldsymbol{G}$ and $\boldsymbol{G}^{\prime}$
- Cyclic Pattern Kernel
- Compare simple cycles in two graphs $\boldsymbol{G}$ and $\boldsymbol{G}^{\prime}$
- Graphlet Kernel
- Count subgraphs of limited size $\boldsymbol{K}$ in $\boldsymbol{G}$ and $\boldsymbol{G}^{\prime}$


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## Shortest-Path Graph Kernel

- Convert graph to all pair shortest path graph
- Floyd-Warshall Algorithm


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## Floyd-Warshall



Original Graph

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## Floyd-Warshall



Original Graph

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## Shortest-Path Graph Kernel

- Apply shortest path kernel
$-K_{s p}\left(G, G^{\prime}\right)=\sum_{e \in E} \sum_{e \prime \in E \prime} K_{\text {walk }}\left(e, e^{\prime}\right)$


## Shortest-Path Graph Kernel

- Apply shortest path kernel
- $K_{s p}\left(G, G^{\prime}\right)=\sum_{e \in E} \sum_{e \prime \in E^{\prime}} K_{\text {walk }}\left(e, e^{\prime}\right)$
- $K_{\text {walk }}\left(e, e^{\prime}\right)=K_{\text {node }}\left(u, u^{\prime}\right) \cdot K_{\text {edge }}\left(e, e^{\prime}\right) \cdot K_{\text {node }}\left(v, v^{\prime}\right)$


## Shortest-Path Graph Kernel

- Apply shortest path kernel
- $K_{s p}\left(G, G^{\prime}\right)=\sum_{e \in E} \sum_{e \prime \in E^{\prime}} K_{\text {walk }}\left(e, e^{\prime}\right)$
- $K_{\text {walk }}\left(e, e^{\prime}\right)=K_{\text {node }}\left(u, u^{\prime}\right) \cdot K_{\text {edge }}\left(e, e^{\prime}\right) \cdot K_{\text {node }}\left(v, v^{\prime}\right)$
- $K_{\text {node }}$ is a valid kernel function for comparing two vertices
- $K_{\text {edge }}$ is a valid kernel function for comparing two edges


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## Shortest Path Graph Kernel

- Lines 4-5 loop through all paths in G1


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## Shortest Path Graph Kernel

- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
$K \leftarrow 0$
$n 1 \leftarrow$ num_node $[g 1]$
$n 2 \leftarrow$ num_node [g2]
for $i=0 \rightarrow n 1, j=0 \rightarrow n 1$ do
if $i \neq j$ AND sp_mat $[g 1][i][j] \neq I N F$ then
for $m=0 \rightarrow n 2, n=0 \rightarrow n 2$ do
if $m \neq n$ AND sp_mat $[g 2][m][n] \neq I N F$ then
$k \_e d g e \leftarrow$ EdgeKernel(sp_mat $[g 1][i][j]$, sp_mat $\left.[g 2][m][n]\right)$
if $K$ _edge $>0$ then
$k \_n o d e 1 \leftarrow \operatorname{NodeKernel}(g 1, g 2, i, m)$
$k \_n o d e 2 \leftarrow \operatorname{NodeKernel}(g 1, g 2, j, n)$
$K+=k \_n o d e 1 * k \_e d g e * k \_n o d e 2$
end if
end if
end for
end if
end for
return $K$
19:


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## Shortest Path Graph Kernel

- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
- Line 8 calculates
$K_{\text {edge }}\left(e, e^{\prime}\right)$

```
\(K \leftarrow 0\)
\(n 1 \leftarrow\) num_node \([g 1]\)
\(n 2 \leftarrow\) num_node [g2]
for \(i=0 \rightarrow n 1, j=0 \rightarrow n 1\) do
    if \(i \neq j\) AND sp_mat \([g 1][i][j] \neq I N F\) then
        for \(m=0 \rightarrow n 2, n=0 \rightarrow n 2\) do
            if \(m \neq n\) AND sp_mat \([g 2][m][n] \neq I N F\) then
            \(k \_e d g e \leftarrow\) EdgeKernel(sp_mat \([g 1][i][j]\), sp_mat \(\left.[g 2][m][n]\right)\)
            if \(K\) _edge \(>0\) then
                    \(k \_n o d e 1 \leftarrow \operatorname{NodeKernel}(g 1, g 2, i, m)\)
                    \(k \_n o d e 2 \leftarrow \operatorname{NodeKernel}(g 1, g 2, j, n)\)
                    \(K+=k \_n o d e 1 * k \_e d g e ~ * k \_n o d e 2\)
            end if
            end if
            end for
        end if
end for
return \(K\)
```

19:

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## Shortest Path Graph Kernel

- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
- Line 8 calculates
$K_{\text {edge }}\left(e, e^{\prime}\right)$
- Lines 10-11 calculate $K_{\text {node }}\left(v, v^{\prime}\right)$
$K \leftarrow 0$
$n 1 \leftarrow$ num_node $[g 1]$
$n 2 \leftarrow$ num_node [g2]
for $i=0 \rightarrow n 1, j=0 \rightarrow n 1$ do
if $i \neq j$ AND sp_mat $[g 1][i][j] \neq I N F$ then for $m=0 \rightarrow n 2, n=0 \rightarrow n 2$ do if $m \neq n$ AND sp_mat $[g 2][m][n] \neq I N F$ then
$k \_$edge $\leftarrow$ EdgeKernel(sp_mat $[g 1][i][j]$, sp_mat $\left.[g 2][m][n]\right)$
if $K$ _edge $>0$ then
$k \_n o d e 1 \leftarrow \operatorname{NodeKernel}(g 1, g 2, i, m)$
$k \_n o d e 2 \leftarrow \operatorname{NodeKernel}(g 1, g 2, j, n)$
$K+=k \_n o d e 1 * k \_e d g e ~ * k \_n o d e 2$
end if
end if
end for
end if
end for
18: return $K$
19:


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## Shortest Path Graph Kernel

- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
- Line 8 calculates
$K_{\text {edge }}\left(e, e^{\prime}\right)$
- Lines 10-11 calculate
$K_{\text {node }}\left(v, v^{\prime}\right)$
- Line 12 calculates
$K_{\text {walk }}\left(e, e^{\prime}\right)$

```
\(K \leftarrow 0\)
\(n 1 \leftarrow\) num_node[g1]
\(n 2 \leftarrow\) num_node[g2]
for \(i=0 \rightarrow n 1, j=0 \rightarrow n 1\) do
    if \(i \neq j\) AND sp_mat \([g 1][i][j] \neq I N F\) then
        for \(m=0 \rightarrow n 2, n=0 \rightarrow n 2\) do
            if \(m \neq n\) AND sp_mat \([g 2][m][n] \neq I N F\) then
            \(k \_e d g e \leftarrow\) EdgeKernel(sp_mat \([g 1][i][j]\), sp_mat \(\left.[g 2][m][n]\right)\)
            if \(K\) _edge \(>0\) then
                    \(k \_n o d e 1 \leftarrow \operatorname{NodeKernel}(g 1, g 2, i, m)\)
                    \(k \_n o d e 2 \leftarrow \operatorname{NodeKernel}(g 1, g 2, j, n)\)
                    \(K+=k \_n o d e 1 * k \_e d g e ~ * k \_n o d e 2\)
            end if
            end if
        end for
        end if
end for
18: return \(K\)
```

19 :

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## Problem to be solved

- Given a set of graphs $\left\{g_{1}, g_{2}, \ldots, g_{n}\right\}$
- Calculate the kernel matrix $\boldsymbol{K}_{n \times n}$
- $K_{(i, j)}$ is the similarity between $g_{i}$ and $g_{j}$


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## GPU Naive1

- Calculate the whole kernel matrix $K_{n \times n}$ in GPU at once
- One GPU thread calculate one element in kernel matrix $K_{n \times n}$
- \# of GPU threads is $n^{2}$


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## Drawbacks of Naive1

- May not have enough GPU memory for large data set
- 3GB global memory on Nvidia Tesla C2050
- GPU threads may have different workload due to different graph sizes
- Load balancing
- Branch divergence
- Works good if all graphs are small and have the same size


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## GPU Naive2

- Calculate similarity between one pair in GPU at a time
- One GPU thread takes one entry in Shortest Path Adjacency Matrix of one input graph
- \# of GPU thread equals to the size of Shortest Path Adjacency Matrix of one input graph
- If there is one edge, then loop through all entries in the other Shortest Path Adjacency Matrix


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## Drawbacks of Naive2

- Waste of GPU resources
- May have idle threads because $\mathbf{0}$ and INF in Shortest Path Adjacency Matrix
- Slow Memory access
- Random, non-coalesced memory access pattern


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## Data Transformation

- Transform Shortest Path Adjacency Matrix to three arrays with length equals to number of shortest paths
- $\boldsymbol{S} \boldsymbol{P}_{-} \boldsymbol{W}$ to store the weight of each path
- SP_S to store the starting node of each path
$-S P_{-} E$ to store the ending node of each path


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| 0 | 1 | 2 |
| :--- | :--- | :--- |
| 0 | 0 | 1 |
| 0 | 0 | 0 |

Shortest Path Adjacency Matrix

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| 0 | 1 | 2 |
| :--- | :--- | :--- |
| 0 | 0 | 1 |
| 0 | 0 | 0 |

Shortest Path Adjacency Matrix


## Advanced GPU Implementation

- Pre-calculation of $K_{\text {node }}$ using Vertex Kernel
- Calculate $K_{\text {walk }}$ uisng Edge Kernel
- Apply Reduction Kernel to sum the results


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Input graphs

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Adjacency matrix

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Adjacency matrix


Shortest Path Adjacency matrix

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| $\zeta_{0}$ | A D | $\zeta_{1}$ | A E | $\zeta_{2}$ | AF |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $\zeta_{3}$ | B D | $\zeta_{4}$ | B E | $\zeta_{5}$ | B F |
| $\zeta_{6}$ | $C D$ | $i_{7}$ | $C E$ | $\zeta_{8}$ | $C F$ |
| Vertex Kernel |  |  |  |  |  |

Adjacency matrix


Shortest Path Adjacency matrix


Input graphs


| $\zeta_{0}$ | A D | $\zeta_{1}$ | A E | $\zeta_{2}$ | AF |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $\zeta_{3}$ | B D | $\zeta_{4}$ | B E | $\zeta_{5}$ | B F |
| $\zeta_{6}$ | $C D$ | $i_{7}$ | $C E$ | $\zeta_{8}$ | $C F$ |
| Vertex Kernel |  |  |  |  |  |

Adjacency matrix


Shortest Path Adjacency matrix

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## Advantage and Disadvantage of Adv. Implementation

- Advantage
- No branch divergence
- Coalesced memory access
- Disadvantage
- Waste of GPU resource when graphs are small


## Hybrid Implementation

- Combine Naive1 and Advanced
- Sort the input graphs according to their sizes
- Set a threshold for the graph size
- For graphs with sizes smaller than threshold, use Naive1
- Otherwise, use Advanced


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## OpenMP Implementation

- Convert the top triangle of the kernel matrix to a 1D array
- Create as many OpenMP threads as number of CPU cores
- Each OpenMP thread calculates one entry in the 1D array in order, goes to next iteration until all entries are computed


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## Execution Environment

- CPU - Intel 5530 Quad core @ 2.4 GHz with 8MB cache (8 OpenMP threads)
- GPU - NVIDIA C2050 (448 Cores @ 1.15GHz) with 3GB GDDR5 1.5 GHZ ECC RAM


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## Synthetic Datasets

| Dataset | Avg. Nodes | Avg. Edges | Avg. SP |
| ---: | ---: | ---: | ---: |
| 10-nodes | 10 | 19 | 61 |
| 20-nodes | 20 | 76 | 367 |
| 30-nodes | 30 | 175 | 867 |
| 40-nodes | 40 | 310 | 1559 |
| 50-nodes | 50 | 489 | 2449 |
| 60-nodes | 60 | 706 | 3540 |
| M1 | 22 | 191 | 930 |
| M2 | 28 | 277 | 1365 |
| M3 | 35 | 362 | 1800 |
| M4 | 41 | 448 | 2235 |
| M5 | 47 | 535 | 2670 |

## Speedups on Uni-size Sets



## Speedups on Mixed Sets



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## Scientific Datasets

| Dataset | Num. of Graphs | Min. Nodes | Max. Nodes | Avg. Nodes | Min. Edges | Max. Edges | Avg. Edges | Avg. SP |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| MUTAG | 188 | 10 | 28 | 18 | 20 | 66 | 39 | 324 |
| ENZYMES | 600 | 2 | 126 | 33 | 2 | 298 | 124 | 1215 |
| NCI1 | 4110 | 3 | 111 | 30 | 4 | 238 | 64 | 1005 |
| NCI109 | 4127 | 4 | 111 | 30 | 6 | 238 | 64 |  |

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## Speedups on Scientific Datasets

| Dataset | Naive 1 | Advanced | Hybrid |
| ---: | ---: | ---: | ---: |
| MUTAG | 2.367 | 1.962 | 2.882 |
| ENZYMES | 1.320 | 10.823 | 10.895 |
| NCI1 | 1.895 | 7.527 | 7.823 |
| NCI109 | 1.992 | 7.751 | 8.037 |

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## Conclusion and Future Work

- We present four different GPU parallelizations
- Achieve up to $44 x$ speedup on synthetic datasets
- Achieve up to $10 x$ speedup on scientific datasets
- We are going to accelerate other graph kernels in the future


## Thanks!

## Questions?

