

# Parallelization of the Shortest Path Graph Kernel on the GPU

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# Outline

- Introduction
  - Graph
  - Graph similarity and graph kernel
  - Shortest Path Graph Kernel
- Parallelization on GPU and CPU
  - Four GPU implementations
  - One OpenMP implementation
- Experiments results
  - Synthetic datasets
  - Scientific datasets
- Conclusion and Future Work



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# Graph

- A graph G is a set of vertices V and edges E, where  $E \subset V^2$
- A graph G may have labels on vertices and/or edges
- The *adjacency matrix* **A** of G is defined as  $[A_{ij}] = \begin{cases} 1 & if(v_i, v_j) \in E\\ 0 & otherwise \end{cases}$



## Labelled Undirected Graphs

vertices





## Labelled Undirected Graphs

vertices

edges







## Labelled Undirected Graphs





## Labelled Undirected Graphs





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# **Graph Similarity**

- How similar are two graphs?
  - For Machine Learning problems like clustering and classification on graphs, graph similarity is crucial.
- Applications
  - Protein function prediction
  - Drug screening
  - Documents classification (Junk mail?)
  - Image classification
  - Cyber security
- Challenges
  - Graph isomorphism is NP-complete
  - Graph comparison via isomorphism is prohibitively expensive



# Graph Kernel

- To Calculate the similarities between two graphs in polynomial time
  - Random Walk Kernel
    - Compare all walks in two graphs **G** and **G'**
  - Shortest Path Kernel
    - Compare all pairs shortest paths for **G** and **G'** via Floyd-Warshall
  - Subtree Kernel
    - Compare subtree-like patterns in two graphs **G** and **G'**
  - Cyclic Pattern Kernel
    - Compare simple cycles in two graphs **G** and **G'**
  - Graphlet Kernel
    - Count subgraphs of limited size K in G and G'



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- Convert graph to all pair shortest path graph
  - Floyd-Warshall Algorithm



## **Floyd-Warshall**





**Original Graph** 

Shortest Path Graph



## **Floyd-Warshall**





**Original Graph** 

Shortest Path Graph



• Apply shortest path kernel

$$- K_{sp(G,G')} = \sum_{e \in E} \sum_{e' \in E'} K_{walk}(e,e')$$



• Apply shortest path kernel

$$- K_{sp(G,G')} = \sum_{e \in E} \sum_{e' \in E'} K_{walk}(e,e') - K_{walk}(e,e') = K_{node}(u,u') \cdot K_{edge}(e,e') \cdot K_{node}(v,v')$$



• Apply shortest path kernel

$$- K_{sp(G,G')} = \sum_{e \in E} \sum_{e' \in E'} K_{walk}(e,e')$$

- $K_{walk}(e, e') = K_{node}(u, u') \cdot K_{edge}(e, e') \cdot K_{node}(v, v')$
- $K_{node}$  is a valid kernel function for comparing two vertices
- $K_{edge}$  is a valid kernel function for comparing two edges



• Lines 4-5 loop through all paths in G1

1:  $K \leftarrow 0$ 2:  $n1 \leftarrow num\_node[g1]$ 3:  $n2 \leftarrow num\_node[g2]$ 4: for  $i = 0 \rightarrow n1$ ,  $j = 0 \rightarrow n1$  do if  $i \neq j$  AND  $sp\_mat[g1][i][j] \neq INF$  then 5:for  $m = 0 \rightarrow n2$ ,  $n = 0 \rightarrow n2$  do 6: if  $m \neq n \ AND \ sp\_mat[g2][m][n] \neq INF$  then 7:  $k\_edge \leftarrow EdgeKernel(sp\_mat[g1][i][j], sp\_mat[g2][m][n])$ 8: if  $K\_edge > 0$  then 9: 10: $k_node1 \leftarrow NodeKernel(q1, q2, i, m)$  $k_node2 \leftarrow NodeKernel(g1, g2, j, n)$ 11: 12: $K + = k\_node1 * k\_edge * k\_node2$ 13:end if end if 14:15:end for 16:end if 17: end for 18: return K19:



- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2

1:  $K \leftarrow 0$ 2:  $n1 \leftarrow num\_node[g1]$ 3:  $n2 \leftarrow num\_node[g2]$ 4: for  $i = 0 \rightarrow n1$ ,  $j = 0 \rightarrow n1$  do if  $i \neq j$  AND  $sp\_mat[g1][i][j] \neq INF$  then 5:for  $m = 0 \rightarrow n2$ ,  $n = 0 \rightarrow n2$  do 6: if  $m \neq n \ AND \ sp\_mat[g2][m][n] \neq INF$  then 7: 8:  $k\_edge \leftarrow EdgeKernel(sp\_mat[g1][i][j], sp\_mat[g2][m][n])$ 9: if  $K\_edge > 0$  then 10: $k_node1 \leftarrow NodeKernel(q1, q2, i, m)$  $k_node2 \leftarrow NodeKernel(g1, g2, j, n)$ 11: 12: $K + = k\_node1 * k\_edge * k\_node2$ 13:end if 14:end if 15:end for 16:end if 17: end for 18: return K19:



- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
- Line 8 calculates
   K<sub>edge</sub>(e, e')

```
1: K \leftarrow 0
 2: n1 \leftarrow num\_node[g1]
 3: n2 \leftarrow num\_node[g2]
 4: for i = 0 \rightarrow n1, j = 0 \rightarrow n1 do
        if i \neq j AND sp\_mat[g1][i][j] \neq INF then
 5:
            for m = 0 \rightarrow n2, n = 0 \rightarrow n2 do
 6:
                if m \neq n \ AND \ sp\_mat[g2][m][n] \neq INF then
 7:
 8:
                     k\_edge \leftarrow EdgeKernel(sp\_mat[g1][i][j], sp\_mat[g2][m][n])
 9:
                     if K\_edge > 0 then
10:
                         k_node1 \leftarrow NodeKernel(g1, g2, i, m)
11:
                         k_node2 \leftarrow NodeKernel(g1, g2, j, n)
12:
                         K + = k\_node1 * k\_edge * k\_node2
13:
                     end if
14:
                 end if
15:
             end for
16:
         end if
17: end for
18: return K
19:
```



- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
- Line 8 calculates
   K<sub>edge</sub>(e, e')
- Lines 10-11 calculate  $K_{node}(v, v')$

1:	$K \leftarrow 0$
2:	$n1 \leftarrow num\_node[g1]$
3:	$n2 \leftarrow num\_node[g2]$
4:	for $i = 0 \rightarrow n1, j = 0 \rightarrow n1$ do
5:	if $i \neq j$ AND $sp\_mat[g1][i][j] \neq INF$ then
6:	for $m = 0 \rightarrow n2$ , $n = 0 \rightarrow n2$ do
7:	if $m \neq n \ AND \ sp\_mat[g2][m][n] \neq INF$ then
8:	$k\_edge \leftarrow EdgeKernel(sp\_mat[g1][i][j], sp\_mat[g2][m][n])$
9:	$ \text{if } K\_edge > 0 \text{ then} \\$
10:	$k\_node1 \leftarrow NodeKernel(g1, g2, i, m)$
11:	$k\_node2 \leftarrow NodeKernel(g1, g2, j, n)$
12:	$K + = k\_node1 * k\_edge * k\_node2$
13:	end if
14:	end if
15:	end for
16:	end if
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- Lines 4-5 loop through all paths in G1
- Lines 6-7 loop through all paths in G2
- Line 8 calculates  $K_{edge}(e, e')$
- Lines 10-11 calculate  $K_{node}(v, v')$
- Line 12 calculates
   K<sub>walk</sub>(e, e')

1:	$K \leftarrow 0$
2:	$n1 \leftarrow num\_node[g1]$
3:	$n2 \leftarrow num\_node[g2]$
4:	for $i = 0 \rightarrow n1$ , $j = 0 \rightarrow n1$ do
5:	if $i \neq j$ AND $sp\_mat[g1][i][j] \neq INF$ then
6:	for $m = 0 \rightarrow n2$ , $n = 0 \rightarrow n2$ do
7:	if $m \neq n \ AND \ sp\_mat[g2][m][n] \neq INF$ then
8:	$k\_edge \leftarrow EdgeKernel(sp\_mat[g1][i][j], sp\_mat[g2][m][n])$
9:	if $K\_edge > 0$ then
10:	$k\_node1 \leftarrow NodeKernel(g1, g2, i, m)$
11:	$k\_node2 \leftarrow NodeKernel(g1, g2, j, n)$
12:	$K + = k\_node1 * k\_edge * k\_node2$
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## Problem to be solved

- Given a set of graphs  $\{g_1, g_2, \dots, g_n\}$
- Calculate the kernel matrix K<sub>nxn</sub>
- $K_{(i,j)}$  is the similarity between  $g_i$  and  $g_j$



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# GPU Naive1

- Calculate the whole kernel matrix  $K_{nxn}$  in GPU at once
  - One GPU thread calculate one element in kernel matrix  $K_{nxn}$
  - # of GPU threads is n<sup>2</sup>



## Drawbacks of Naive1

- May not have enough GPU memory for large data set
   3GB global memory on Nvidia Tesla C2050
- GPU threads may have different workload due to different graph sizes
  - Load balancing
  - Branch divergence
- Works good if all graphs are small and have the same size



# GPU Naive2

- Calculate similarity between one pair in GPU at a time
  - One GPU thread takes one entry in Shortest Path
     Adjacency Matrix of one input graph
  - # of GPU thread equals to the size of Shortest Path
     Adjacency Matrix of one input graph
  - If there is one edge, then loop through all entries in the other Shortest Path Adjacency Matrix



## Drawbacks of Naive2

- Waste of GPU resources
  - May have idle threads because *0* and *INF* in Shortest Path Adjacency Matrix
- Slow Memory access
  - Random, non-coalesced memory access pattern



#### **Shortest Path Adjacency Matrix**



## Data Transformation

- Transform Shortest Path Adjacency Matrix to three arrays with length equals to number of shortest paths
  - **SP\_W** to store the weight of each path
  - SP\_S to store the starting node of each path
  - **SP\_E** to store the ending node of each path





Shortest Path Adjacency Matrix





#### Shortest Path Adjacency Matrix





## Advanced GPU Implementation

- Pre-calculation of K<sub>node</sub> using Vertex Kernel
- Calculate K<sub>walk</sub> uisng Edge Kernel
- Apply *Reduction Kernel* to sum the results





Input graphs





Input graphs

	Α	в	С		D	E	F
А	0	1	0	D	0	1	1
В	0	0	1	E	0	0	0
С	0	0	0	F	0	0	0

Adjacency matrix





Input graphs

	Α	В	С		D	E	F	
А	0	1	0	D	0	1	1	
В	0	0	1	E	0	0	0	
С	0	0	0	F	0	0	0	

Adjacency matrix



Shortest Path Adjacency matrix







Shortest Path Adjacency matrix

30 A D	, A Ε	3,2 A F
≩ <sub>3</sub> βD	<b>≩</b> 4 ВЕ	, ≩5 BF
3 <sub>6</sub> ⊂ D	, ₹7 СЕ	3,8 CF

Vertex Kernel





Shortest Path Adjacency matrix

3 <sub>0</sub> ad	<li>↓ A E</li>	3,2 A F
≩ <sub>3</sub> ΒD	<b>≩</b> 4 ВЕ	3,5 BF
3 <sub>6</sub> ⊂ D	ζ <sub>7</sub> сε	3,8 CF

Vertex Kernel



Edge Kernel



# Advantage and Disadvantage of Adv. Implementation

- Advantage
  - No branch divergence
  - Coalesced memory access
- Disadvantage
  - Waste of GPU resource when graphs are small



## Hybrid Implementation

- Combine Naive1 and Advanced
- Sort the input graphs according to their sizes
- Set a threshold for the graph size
  - For graphs with sizes smaller than threshold, use
     *Naive1*
  - Otherwise, use *Advanced*



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## **OpenMP** Implementation

- Convert the top triangle of the kernel matrix to a 1D array
- Create as many OpenMP threads as number of CPU cores
- Each OpenMP thread calculates one entry in the 1D array in order, goes to next iteration until all entries are computed



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## **Execution Environment**

- **CPU** Intel 5530 Quad core @ 2.4 GHz with 8MB cache (8 OpenMP threads)
- **GPU** NVIDIA C2050 (448 Cores @ 1.15GHz) with 3GB GDDR5 1.5 GHZ ECC RAM



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#### Synthetic Datasets

Dataset	Avg. Nodes	Avg. Edges	Avg. SP
10-nodes	10	19	61
20-nodes	20	76	367
30-nodes	30	175	867
40-nodes	40	310	1559
50-nodes	50	489	2449
60-nodes	60	706	3540
M1	22	191	930
M2	28	277	1365
M3	35	362	1800
M4	41	448	2235
M5	47	535	2670



#### Speedups on Uni-size Sets





#### Speedups on Mixed Sets





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#### **Scientific Datasets**

Dataset	Num. of Graphs	Min. Nodes	Max. Nodes	Avg. Nodes	Min. Edges	Max. Edges	Avg. Edges	Avg. SP
MUTAG	188	10	28	18	20	66	39	324
ENZYMES	600	2	126	33	2	298	124	1215
NCI1	4110	3	111	30	4	238	64	1005
NCI109	4127	4	111	30	6	238	64	995



## Speedups on Scientific Datasets

Dataset	Naive1	Advanced	Hybrid
MUTAG	2.367	1.962	2.882
ENZYMES	1.320	10.823	10.895
NCI1	1.895	7.527	7.823
NCI109	1.992	7.751	8.037



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# **Conclusion and Future Work**

- We present four different GPU parallelizations
- Achieve up to 44x speedup on synthetic datasets
- Achieve up to 10x speedup on scientific datasets
- We are going to accelerate other graph kernels in the future



# Thanks! Questions?